# CS856: Presentation of "Rollback-Recovery for Middleboxes"

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## Motivation

- Dedicated middlebox hardware, with backup, is expensive
- Middlebox applications that use NFV run on diverse hardware platforms, which have a higher probability of failure

# Middlebox application model



#### Challenges to recovery

- Statefulness: Shared variables (e.g. counters) need to be restored before processing new packets
- Non-determinism: Access order to shared variables is important; access of hardware clocks needs to be reproduced for stateful recovery
- Low latency: Normal operation needs to be in the order of microseconds

## Replay vs No-Replay designs

- No-Replay: Snapshot of system and buffering of output until next snapshot. Simple, but slow.
- Replay: Snapshot of system and write-ahead-logging of input in between snapshots. Output is only released after input is safely logged. Lower latency per operation, but more expensive logging.

# FTMB architecture



## Packet dependencies



Packet can be released when all its causal dependencies have been recorded to safe storage as Packet Access Logs (PALs).

# Parallel release

- Master: Packet is released together with a vector clock representing the number of PALs each queue has processed: e.g. next packet has vector clock [56, 77, 63, 77].
- Output Logger: Packet is released when each queue has processed more PALs than the current value of the vector clock: e.g. [45, 76, 60, 70] can't be released because of the third queue.





#### Implementation

- Input Logger: Buffers incoming packets, in between snapshots, for replay. Adds one hop delay.
- Master: Processes packet. Sends PALs and output packets to the stable storage and output logger. Adds processing + PAL generation and transmission delay.
- Output Logger: Buffers outgoing packets until the vector clocks increment to the per-packet vector clock value.
  Generally adds the one hop latency to assert that PALs have been stored.

# Discussion

- Middlebox application code has to be annotated. Modulo that, the solution is generic.
- Performance numbers show feasibility of approach.
- The Input Logger, Stable Storage and Output Logger are now the points of failure.
- Causal consistency work in databases has more depth, and covers the "novelty" presented here. The application of causal consistency to middlebox code, however, might be novel.

# Causal consistency work

- Don't Settle for Eventual: Scalable Causal Consistency for Wide-area Storage with COPS
- Stronger semantics for low-latency geo-replicated storage
- Orbe: scalable causal consistency using dependency matrices and physical clocks